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# Complete catalogue of graphs of maximum degree 3 and defect at most 4

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#### ABSTRACT

We consider graphs of maximum degree 3, diameter  $D \ge 2$  and at most 4 vertices less than the Moore bound  $M_{3,D}$ , that is,  $(3,D,-\epsilon)$ -graphs for  $\epsilon \le 4$ .

We prove the non-existence of (3, D, -4)-graphs for  $D \ge 5$ , completing in this way the catalogue of  $(3, D, -\epsilon)$ -graphs with  $D \ge 2$  and  $\epsilon \le 4$ . Our results also give an improvement to the upper bound on the largest possible number  $N_{3,D}$  of vertices in a graph of maximum degree 3 and diameter D, so that  $N_{3,D} \le M_{3,D} - 6$  for  $D \ge 5$ .

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#### 1. Introduction

The optimality of a network has been interpreted in various ways, see, for instance, [16]. One possible interpretation can be stated as follows:

An *optimal network* contains the maximum possible number of nodes, given a limitation on the number of connections attached to a node and a limitation on the number of traversed links between any two farthest nodes.

In graph-theoretical terms, the preceding interpretation leads to the

*Degree/diameter problem*: Given natural numbers  $\Delta \geq 2$  and  $D \geq 1$ , find the largest possible number of vertices  $N_{\Delta,D}$  in a graph of maximum degree  $\Delta$  and diameter D.

It is straightforward to verify that  $N_{\Delta,D}$  is defined for  $\Delta \geq 2$  and  $D \geq 1$ . An upper bound on  $N_{\Delta,D}$  is given by the following expression [3,13].

$$N_{\Delta,D} \leq 1 + \Delta + \Delta(\Delta - 1) + \dots + \Delta(\Delta - 1)^{D-1}$$

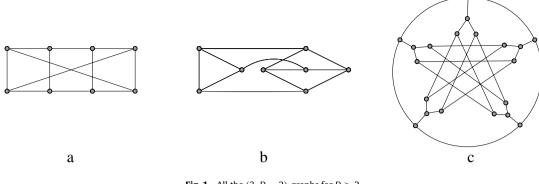
$$= 1 + \Delta[1 + (\Delta - 1) + \dots + (\Delta - 1)^{D-1}]$$

$$= \begin{cases} 1 + \Delta \frac{(\Delta - 1)^{D} - 1}{\Delta - 2} & \text{if } \Delta > 2\\ 2D + 1 & \text{if } \Delta = 2. \end{cases}$$
(1)

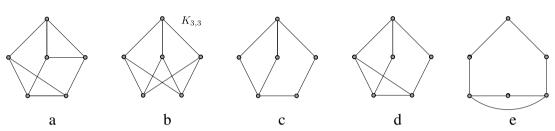
This expression is known as the *Moore bound*, and is denoted by  $M_{\Delta,D}$ . A graph whose order is equal to the Moore bound is called a *Moore graph*.

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**Fig. 1.** All the (3, D, -2)-graphs for  $D \ge 2$ .



**Fig. 2.** All the (3, 2, -4)-graphs.

Moore graphs exist only for certain special values of maximum degree and diameter. To be more precise, for diameter D=1 and degree  $\Delta\geq 1$ , Moore graphs are the complete graphs of order  $\Delta+1$ . For diameter D=2 Hoffman and Singleton [9] proved that Moore graphs exist for  $\Delta=2,3,7$  and possibly 57, but not for any other degree. Finally, for  $D\geq 3$  and  $\Delta=2$ , Moore graphs are the cycles on 2D+1 vertices. The fact that Moore graphs do not exist for  $D\geq 3$  and  $\Delta\geq 3$  was shown by Damerell [5] and, independently, also by Bannai and Ito [1].

Therefore, we are interested in studying the existence of large graphs of given maximum degree  $\Delta$ , diameter D and order  $M_{\Delta,D} - \epsilon$ , for  $\epsilon > 0$ , that is,  $(\Delta, D, -\epsilon)$ -graphs, where  $\epsilon$  is called the *defect*.

Since the case  $\Delta=2$  is completely settled  $(N_{2,D}=2D+1, \text{ for }D\geq 3)$ , in this paper, we consider the next case,  $\Delta=3$ . For  $D\geq 2$ , if a  $(3,D,-\epsilon)$ -graph had a vertex of degree at most 2 then the order of such a graph would be at most  $\frac{2}{3}M_{3,D}+\frac{1}{3}$ ; see [10]. Therefore, we can state the following proposition.

**Proposition 1.1** ([10]). If  $\epsilon < \frac{M_{3,D}}{3} - \frac{1}{3}$  then a  $(3, D, -\epsilon)$ -graph is regular.

By Proposition 1.1, for  $\epsilon < \frac{M_{3,D}}{3} - \frac{1}{3}$ , odd  $\epsilon$ , and  $D \ge 2$ , a  $(3,D,-\epsilon)$ -graph is *cubic*, and must have an even number of vertices. Therefore, these graphs do not exist when  $\epsilon = 1,3$ . Thus, the next interesting cases occur when  $\epsilon = 2$  and 4.

The case of  $\epsilon=2$  was analyzed by Jørgensen [10]. Jørgensen proved that for  $D\geq 4$  there are no (3, D, -2)-graphs and showed the uniqueness of the two known (3, 2, -2)-graphs (graphs (a) and (b) in Fig. 1) and of the (3, 3, -2)-graph (graph (c) in Fig. 1).

The case  $\epsilon=4$  and D=2 or 3 was considered in [14], where we presented all the (3,2,-4)-graphs. The unique (3,3,-4)-graph was constructed initially by Faradžev [8], and later rediscovered by McKay and Royle [12], who proved its uniqueness; see Figs. 2 and 3.

For diameter 4 the non-existence of (3, 4, -4)-graphs was proved by Jørgensen [11].

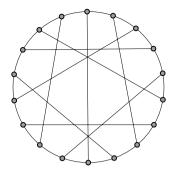
A simple counting argument shows that a (3, D, -4)-graph,  $D \ge 3$ , has girth at least 2D - 2. In [14] we proved that the girth must be at least 2D - 1, and conjectured that its real value is 2D.

In this paper we prove that if a (3, D, -4)-graph with  $D \ge 5$  exists then its girth must be 2D. Moreover, using this result about the girth of such graphs, we show that there are no (3, D, -4)-graphs for  $D \ge 5$ , thus completing the census of  $(3, D, -\epsilon)$ -graphs with  $D \ge 2$  and  $\epsilon \le 4$ .

Note that some parts of our proof are inspired by the reasoning used by Jørgensen in [11].

Values of  $N_{3,D}$  are known only for D=2, 3 and 4. For D=2,  $N_{3,2}=M_{3,2}$ , and the unique graph is the Petersen graph; see [9]. For D=3,  $N_{3,3}=M_{3,3}-2$ , and the unique graph, depicted in Fig. 1(c), was found by Bermond, Delorme and Farhi [2,10]. For D=4, by proving the non-existence of (3, 4, -6)-graphs, Buset[4] showed that  $N_{3,4}=M_{3,4}-8$ , and the two known (non-isomorphic) graphs, constructed by Doty [7] and by von Conta [15], therefore became the largest graphs when  $\Delta=3$  and D=4.

Our results give an improvement on the upper bound of  $N_{3,D}$ , so that  $N_{3,D} \leq M_{3,D} - 6$  for  $D \geq 5$ .



**Fig. 3.** The unique (3, 3, -4)-graph.

The rest of this paper is structured as follows: in Section 2, we settle the notation and terminology used throughout this paper and we give some preliminary results. Section 3 is devoted to proving that if a (3, D, -4)-graph with  $D \ge 5$  exists then it must have girth 2D. In Section 4 we prove the non-existence of (3, D, -4)-graphs with  $D \ge 5$ ; and in Section 5 we give a summary of our results.

It is perhaps worth noting that the case of (3, D, -4)-graphs is particularly interesting, because it is the first result concerning  $(\Delta, D, -\epsilon)$ -graphs of defect greater than the maximum degree of the graph.

### 2. Terminology and preliminary results

All graphs considered in this paper are simple, that is, they have neither loops nor multiple edges.

Throughout this paper, it is assumed that the reader is already familiar with basic graph theory, and therefore with its main concepts and results. Thus, the only objective of this section is to settle those notations that could vary among texts.

The terminology and notation used in this paper is standard and consistent with that used in [6].

The vertex set of a graph  $\Gamma$  is denoted by  $V(\Gamma)$ , and its edge set by  $E(\Gamma)$ . In  $\Gamma$  a vertex of degree at least 3 is called a branch vertex of  $\Gamma$ . For an edge  $e = \{x, y\}$ , we write e = xy, or simply xy, or alternatively,  $x \sim y$ . If two vertices u and u are not adjacent then we write  $x \nsim y$ . The length of a path P is the number of edges in P. A path of length k is called a k-path. A path from a vertex x to a vertex y is denoted by x - y. Whenever we refer to paths, we mean shortest paths. A cycle of length k is called a k-cycle.

We will also use the following notations for subpaths of a path  $P = x_0 x_1 \dots x_k : x_i P x_j = x_i \dots x_j$ , where  $0 \le i \le j \le k$ .

The set of vertices at distance k from a vertex x is denoted by  $N_k(x)$ . The set of neighbors of a vertex x in  $\Gamma$  is simply denoted by N(x). The set of edges in the graph  $\Gamma$  joining a vertex x in  $X \subseteq V(\Gamma)$  to a vertex y in  $Y \subseteq V(\Gamma)$  is denoted by E(X,Y); for simplicity, instead of E(X,X), we write E(X).

The *difference* between the graphs  $\Gamma$  and  $\Gamma'$ , denoted by  $\Gamma - \Gamma'$ , is the graph with vertex set  $V(\Gamma) - V(\Gamma')$  and edge set formed by all the edges with both endvertices in  $V(\Gamma) - V(\Gamma')$ .

The union of three independent paths of length D with common endvertices is denoted by  $\Theta_D$ .

Finally, we call a cycle of length at most 2D a short cycle, and we call a vertex x a saturated vertex if x cannot belong to any further short cycle.

From now on, let  $\Gamma$  be a (3, D, -4)-graph for  $D \ge 5$ . By Proposition 1.1,  $\Gamma$  must be regular.

Furthermore, we have

**Proposition 2.1** ([14]). A (3, D, -4)-graph for  $D \ge 5$  has girth at least 2D - 1.

In [14] it was conjectured that

**Conjecture 2.1** ([14]). The girth of a (3, D, -4)-graph,  $D \ge 5$ , is 2D.

If the girth of  $\Gamma$  is 2D-1 then there exists a vertex x in  $\Gamma$  such that x lies on either one or two (2D-1)-cycles. Note that no vertex x can lie on more than two such cycles, otherwise  $|E(N_{D-1}(x))| \ge 3$ , implying  $\Gamma \le M_{3,D}-6$ , a contradiction.

Using a simple counting argument, we classify each vertex of a (3, D, -4)-graph according to the short cycles on which the vertex lies, as shown in Proposition 2.2.

**Proposition 2.2.** Let x be a vertex of  $\Gamma$ . Then x lies on the short cycles specified below, and no other short cycle. We have the following cases:

*x* is contained in two (2D-1)- cycles. Then

(i) x **lies on exactly two** (2D-1)- **cycles** whose intersection is an l-path for some l such that  $1 \le l \le D-1$ . If l = D-1 then x is also contained in one 2D-cycle; or

x is contained in exactly one (2D-1)- cycle. Then also

- (ii) x is a branch vertex of one  $\Theta_D$ , or
- (iii) x is contained in exactly two 2D-cycles; or

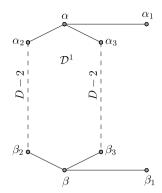


Fig. 4. Auxiliary figure for Lemma 2.1.

### x is contained in no (2D-1)- cycle. Then also

- (iv) x is a branch vertex of exactly two  $\Theta_D$ , or
- (v) x is a branch vertex of one  $\Theta_D$ , and is contained in two more 2D-cycles, or
- (vi) *x* is contained in exactly four 2D-cycles.

Each case is considered as a type. For instance, a vertex satisfying case (i) is called a vertex of Type (i).

**Proof.** By Proposition 2.1, we see that  $N_i(x)$  is an independent set, for  $i \in \{1, ..., D-2\}$ , and that  $|N_{D-1}(x)| = 3 \times 2^{D-2}$ . It is clear that  $|E(N_{D-1}(x))| \le 2$ , otherwise  $|\Gamma| \le M_{3,D} - 6$ . We distinguish three cases according to the possible values of  $|E(N_{D-1}(x))|$ .

Case 1.  $|E(N_{D-1}(x))| = 2$ .

In this case these two edges either have a common endvertex or are independent. Therefore, (i) follows.

Case 2.  $|E(N_{D-1}(x))| = 1$ . Since  $|N_D(x)| = 3 \times 2^{D-1} - 4$  and  $|E(N_{D-1}(x), N_D(x))| = 3 \times 2^{D-1} - 2$ , we obtain (ii) or (iii).

Case 3.  $|E(N_{D-1}(x))| = 0$ . Since  $|N_D(x)| = 3 \times 2^{D-1} - 4$  and  $|E(N_{D-1}(x), N_D(x))| = 3 \times 2^{D-1}$ , it follows that x is a vertex of Type (iv), (v) or (vi).  $\square$ 

**Observation 2.1.** If a vertex  $x \in \Gamma$  belongs to exactly one (2D-1)-cycle  $C^1$  then the intersection of  $C^1$  and any 2D-cycle is a path of length at most D-1.

Next we prove a lemma that will be used repeatedly in the rest of this paper.

**Lemma 2.1** (Intersection Lemma). Let  $\mathcal{D}^1$  be a 2D-cycle in  $\Gamma$ . Let  $\alpha$  and  $\beta$  be vertices on  $\mathcal{D}^1$  such that  $d(\alpha, \beta) = D$ . Let  $\alpha_1$  be the neighbor of  $\alpha$  not contained in  $\mathcal{D}^1$ . Let us suppose that  $\alpha$  is not a branch vertex of a  $\Theta_D$ , and that  $\alpha_1$  is contained in at most one (2D-1)-cycle, say  $\mathcal{C}$ , which also contains  $\alpha$ . Then

- (i) the intersection of  $\mathcal{D}^1$  and  $\mathcal{C}$  is a path of length D-1, or
- (ii) there exists another 2D-cycle, say  $\mathcal{D}^2$ , containing  $\alpha$  and  $\alpha_1$ . Furthermore, the intersection of  $\mathcal{D}^1$  and  $\mathcal{D}^2$  is a path of length D-1.

**Proof.** Let  $\mathcal{D}^1$  be a 2*D*-cycle of  $\Gamma$ , and let  $\alpha$ ,  $\alpha_1$ ,  $\alpha_2$ ,  $\alpha_3$ ,  $\beta$ ,  $\beta_1$ ,  $\beta_2$  and  $\beta_3$  be as in Fig. 4.

Let  $P^1 = \alpha_1 - \beta$ . The length of  $P^1$  must be D, since  $\alpha$  is not a branch vertex of a  $\Theta_D$ ,  $\alpha_1$  is contained in at most one (2D-1)-cycle, and the girth of  $\Gamma$  is at least 2D-1. Therefore, we have two possibilities: either  $P^1$  goes through  $\beta_2$  or  $\beta_3$ , or it goes through  $\beta_1$ . In the first case  $V(P^1 \cap \mathcal{D}^1) = \{\beta_2 \text{ or } \beta_3, \beta\}$ , and (i) follows. In the second case we consider the neighbor  $\alpha'$  of  $\alpha_1$  such that  $\alpha' \neq \alpha$  and  $\alpha' \notin P^1$ . A path  $P^2 = \alpha' - \beta$  does not pass through  $\beta_1$ , otherwise  $\alpha_1$  would belong to a cycle of length at most 2D-1 that does not contain  $\alpha$ , contradicting our assumptions. Therefore,  $P^2$  is a path of length either D-1 or D, which goes through  $\beta_2$  or  $\beta_3$ , and  $V(P^2 \cap \mathcal{D}^1) = \{\beta_2 \text{ or } \beta_3, \beta\}$ . Consequently, if  $P^2$  is a (D-1)-path then (i) follows, otherwise (ii) follows.

Note that if  $\alpha_1$  is contained in no (2D-1)-cycle then (ii) follows.

## 3. On the girth of (3, D, -4)-graphs with $D \ge 5$

The aim of this section is to prove that the girth of  $\Gamma$  is exactly 2D. This result will be obtained by ruling out the existence of vertices of Type (i), (ii) or (iii).

**Theorem 3.1.** A(3, D, -4)-graph,  $D \ge 5$ , does not contain a vertex of Type (i), (ii) or (iii).

We prove Theorem 3.1 by eliminating, in order, the existence of vertices of each type under consideration.

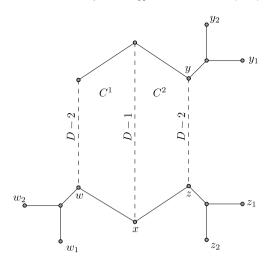


Fig. 5. Auxiliary figure for Lemma 3.1.

Non-existence of vertices of Type (i)

In the next two lemmas we give some necessary conditions for the existence of vertices of Type (i).

**Lemma 3.1.** Let x be a vertex lying on two (2D-1)-cycles. Then the intersection of such (2D-1)-cycles is a path of length at most D-2.

**Proof.** We proceed by way of contradiction. Let us consider a vertex  $x \in \Gamma$  lying on two (2D-1)-cycles, say  $C^1$  and  $C^2$ , and let us further suppose that the intersection of  $C^1$  and  $C^2$  is a path of length D-1. Then  $\Gamma$  contains the subgraph in Fig. 5.

Let  $x, w, w_1, w_2, z, z_1, z_2, y, y_1, y_2$  be as in Fig. 5. A path  $P = y_1 - x$  is a D-path, since  $d(x, y_1) \le D$ , and by Proposition 2.2(i). Besides, by Proposition 2.2(i), if P intersects with  $C^1$  then  $V(P \cap C^1) = \{x, w\}$ ; if instead P intersects with  $C^2$  then  $V(P \cap C^2) = \{x, z\}$ . Therefore, P should pass through either  $w_1, w_2, z_1$ , or  $z_2$ .

If the path  $P = y_1 - x$  passed through either  $z_1$  or  $z_2$ , say  $z_1$ , then z would be contained in two (2D - 1)-cycles, namely,  $C^2$  and  $C = zC^2y \dots y_1Pz_1 \dots z$ , and in one 2D-cycle. However, by Proposition 2.2(i), the intersection of C and  $C^2$  should be a path of length D - 1, and in this case, the intersection is a path of length D - 2, namely,  $zC^2y$ , a contradiction. Therefore,  $P = y_1 - x$  passes through either  $w_1$  or  $w_2$ . Analogously, a path  $Q = y_2 - x$  is a D-path, and passes through either  $w_1$  or  $w_2$ . Thus,  $\Gamma$  contains a cycle of length at most 2D - 2, contradicting Proposition 2.1.  $\square$ 

**Lemma 3.2.** Let us assume that  $\Gamma$  contains two non-disjoint (2D-1)-cycles. Then the intersection of such (2D-1)-cycles is a path of length exactly D-2.

**Proof.** Let us suppose that  $\Gamma$  contains two non-disjoint (2D-1)-cycles, denoted by  $C^1$  and  $C^2$ .

To prove this lemma we proceed by way of contradiction. Suppose that the intersection of the cycles  $C^1$  and  $C^2$  is a path of length l, with  $l \in \{1, \ldots, D-3\}$ . Recall that the case of l = D-1 is ruled out by Lemma 3.1.

As  $C^1 \neq C^2$ , there are two vertices x and  $x_1$  such that  $x \in (C^1 \cap C^2)$ ,  $x_1 \in (C^2 - C^1)$  and  $x \sim x_1$ . We may also assume that  $x_1$  has a neighbor  $x_3$  such that  $x \neq x_3$  and  $x_3 \in C^2$ .

Let  $x_4$ , y,  $y_1$ ,  $y_2$ ,  $y_3$ ,  $y_4$ , z,  $z_1$ ,  $z_2$ ,  $z_3$ , and  $z_4$  be as in Fig. 6(a).

Let us first consider a path  $P^1 = x_3 - z$ . Since the intersection of  $C^1$  and  $C^2$  is a path of length l with  $1 \le l \le D - 3$ , we have  $x_1 \notin P^1$ . By assumption,  $P^1$  cannot go through  $z_1$ .  $P^1$  does not pass through y, and  $V(P^1 \cap C^1) = \{z\}$ , since x is a vertex of Type (i), and the intersection of  $C^1$  and  $C^2$  is a path of length l with  $l \in \{1, \ldots, D-3\}$ . Therefore,  $P^1$  is a D-path that passes through either  $z_3$  or  $z_4$ , say  $z_3$ . By following similar reasoning, a path  $P^2 = x_4 - z$  goes through either  $z_3$  or  $z_4$ . Then  $V(P^1 \cap P^2) = \{z_2, z\}$ , otherwise there would be a cycle of length at most (2D-2). Therefore, the path  $P^2$  uses the vertex  $z_4$ , and is a D-path. Note that  $x_1x_3P^1z_3z_2z_4P^2x_4x_1$  is a 2D-cycle, denoted by  $D^1$ .

In the same way, we can assume that the paths  $Q^1 = x_3 - y$  and  $Q^2 = x_4 - y$  use the vertices  $y_3$  and  $y_4$ , respectively. Furthermore, both  $Q^1$  and  $Q^2$  are D-paths, and  $V(Q^1 \cap Q^2) = \{y_2, y\}$ . Note that  $x_1x_3Q^1y_3y_2y_4Q^2x_4x_1$  is also a 2D-cycle, denoted by  $D^2$ . Thus,  $x_1$  and  $x_3$  are contained in the 2D-cycles  $D^1$  and  $D^2$ , and in the (2D-1)-cycle  $C^2$ .

Let *s* and *r* be the neighbors of  $x_3$  different from  $x_1$  such that  $s \in P^1$  and  $r \notin P^1$ .

Recall that the vertices  $x_1$  and  $x_3$  cannot be contained in any additional cycle of length at most 2D.

We distinguish two cases: either  $V(P^1 \cap Q^1) = \{x_3, s\}$  or  $V(P^1 \cap Q^1) = \{x_3\}$ .

Case 1. The paths  $P^1$  and  $Q^1$  intersect at  $x_3$  and s.

In this case the vertices y and z lie on a further (2D-1)-cycle, namely,  $sQ^1y_3y_2yzz_2z_3P^1s$ , and consequently, the paths  $P^2$  and  $Q^2$  should intersect only at  $x_4$ , otherwise y and z would be contained in three (2D-1)-cycles; see Fig. 6(b).

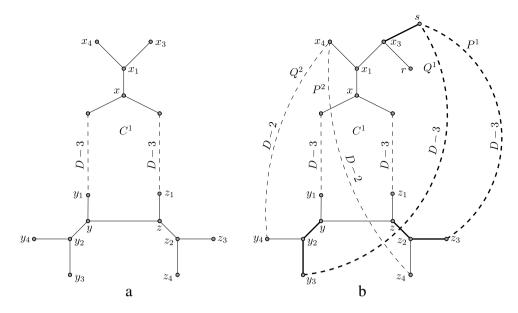


Fig. 6. (a): Auxiliary figure for Lemma 3.2. (b): Auxiliary figure for Case 1 of Lemma 3.2.

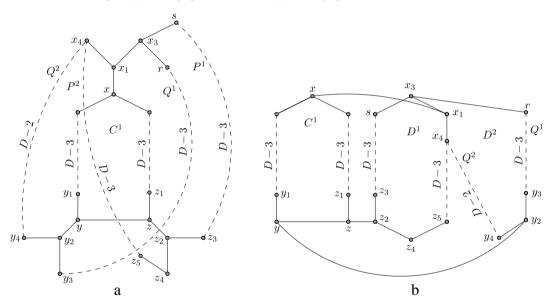


Fig. 7. Auxiliary figure for Case 2 of Lemma 3.2.

Let us consider a path A = r - z. Considering our assumptions, and that the vertices  $x_1$  and  $x_3$  cannot belong to a further short cycle, we have that the path A cannot use the vertices  $z_1$ ,  $z_3$ ,  $z_4$ ,  $y_1$ ,  $y_3$  or  $y_4$ , and therefore, r cannot reach z in at most D steps, a contradiction.

Case 2. The paths  $P^1$  and  $Q^1$  intersect only at  $x_3$ .

We may assume that the paths  $P^2$  and  $Q^2$  intersect only at  $x_4$ ; see Fig. 7(a).

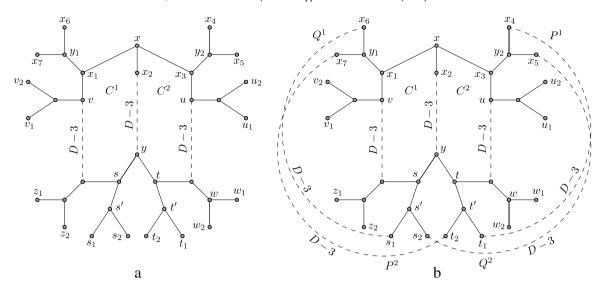
Without loss of generality, as  $x_3 \in C^2$ , we may also assume that  $r \in C^2$ .

Let  $z_5$  be as in Fig. 7(a). Note that since  $D \ge 5$ ,  $\Gamma$  contains the subgraph depicted in Fig. 7(a).

To achieve a better understanding of this case, we depict Fig. 7(a) in a different way, by drawing our attention to the vertices  $x_3$  and r, and to the 2*D*-cycle  $D^1$ ; see Fig. 7(b).

We see that the premises of the Intersection Lemma hold. Mapping the vertex  $x_3$  to  $\alpha$ , r to  $\alpha_1$ ,  $z_4$  to  $\beta$ , and mapping the 2*D*-cycle  $D^1$  to  $D^1$ , and the (2D-1)-cycle  $C^2$  to C, we obtain, by the Intersection Lemma, that one of the following cases holds. In the first case,  $x_3$  and r are contained in a (2D-1)-cycle that intersects with  $D^1$  at a path of length D-1. This cycle would be precisely  $C^2$ , implying D-1=1, a contradiction. In the second case,  $x_3$  and r are contained in another 2*D*-cycle that intersects with  $D^1$  at a path of length D-1. This cycle would be precisely  $D^2$ , implying D-1=2, a contradiction.

Consequently, the cycles  $C^1$  and  $C^2$  intersect at a path of length exactly D-2.



**Fig. 8.** Vertex of Type (i) in a (3, D, 4)-graph when  $D \ge 5$ .

Using the structural results from Lemmas 3.1 and 3.2, the next proposition rules out the existence of vertices of Type (i).

**Proposition 3.1.** A(3, D, -4)-graph, D > 5, does not contain a vertex of Type (i).

**Proof.** Let  $x \in \Gamma$  be a vertex of Type (i), lying on the (2D-1)-cycles  $C^1$  and  $C^2$ . In view of Lemma 3.2, the intersection of  $C^1$  and  $C^2$  is a path of length D-2. Then since  $D \ge 5$ ,  $\Gamma$  contains the subgraph in Fig. 8(a).

Let x,  $x_1$ ,  $x_2$ ,  $x_3$ ,  $x_4$ ,  $x_5$ ,  $x_6$ ,  $x_7$ , y,  $y_1$ ,  $y_2$ , u,  $u_1$ ,  $u_2$ , v,  $v_1$ ,  $v_2$ , s, s',  $s_1$ ,  $s_2$ , t, t',  $t_1$ ,  $t_2$ ,  $t_1$ ,  $t_2$ ,  $t_3$ ,  $t_4$ , and  $t_5$  be as in Fig. 8(a). Let us first consider a path  $t^2 = t_1 - t_2$ . Note that since all the vertices in the intersection of  $t^2 = t_1 - t_2$ . follows that t',  $x_2 \notin P^1$ . The vertex u is not contained in  $P^1$ , otherwise there would exist a cycle of length at most 2D-2. The vertex v is not contained in  $P^1$  either, otherwise the vertex y would be contained in two (2D-1)-cycles  $(C^1$  and  $C^2)$ , and in one further cycle of length at most 2D, a contradiction. Therefore,  $P^1$  is a D-path, and passes through either  $x_4$ ,  $x_5$ ,  $x_6$ or  $x_7$ . Analogously, a path  $P^2 = t_2 - x$  is a D-path, and goes through either  $x_4, x_5, x_6$  or  $x_7$ . Moreover, neither  $x_1 \in (P^1 \cap P^2)$ nor  $x_3 \in (P^1 \cap P^2)$ , otherwise there would exist a cycle of length at most 2D-2 in  $\Gamma$ . Without loss of generality, we may therefore assume that  $P^1$  goes through  $x_4$ , and that  $P^2$  goes through  $x_7$ .

In the same way, we may assume that the D-paths  $Q^1 = s_1 - x$  and  $Q^2 = s_2 - x$  go through  $x_6$  and  $x_5$ , respectively. Note that  $V(P^2 \cap Q^1) = \{y_1, x_1, x\}$ , otherwise y would be contained in an additional cycle of length at most 2D. Analogously,  $V(P^1 \cap Q^2) = \{y_2, x_3, x\}$ ; see Fig. 8(b).

Note also that  $D^1 = tt't_1P^1x_4y_2x_3uC^2t$  and  $D^2 = ss's_1Q^1x_6y_1x_1vC^1s$  are 2D-cycles.

Let  $x_8$  and  $x_9$  be the neighbors of  $x_6$  different from  $y_1$  such that  $x_8 \in Q^1$  and  $x_9 \notin Q^1$ . Let r be the neighbor of v on  $C^1$ different from  $x_1$ ; see Fig. 9(a).

The paths  $M^1 = w_1 - x$  and  $M^2 = w_2 - x$  are D-paths. Note that  $w \notin M^1$ . The vertex  $u \notin M^1$ , otherwise there would exist a cycle of length at most 2D-3 in  $\Gamma$ . The vertex  $x_2 \notin M^1$ , otherwise  $x_2$  would be contained in a further cycle of length at most 2D, contradicting Proposition 2.2(i). Since the vertex  $x_3$  is of Type (iii) or (iv) (see Proposition 2.2(iii) and (iv)),  $x_3 \notin M^1$ , otherwise  $x_3$  would be contained in an additional cycle of length at most 2D-1 ( $x_3$  is already contained in the (2D-1)-cycle  $C^2$  and in the 2D-cycle  $D^1$ ). Furthermore, if the path  $M^1$  went through  $x_7$  then the vertex t would be contained in an additional cycle of length at most 2D-1, contradicting the fact that  $t \in \text{Type (iii)}$  or (iv)  $(t \in C^2 \text{ and } D^1)$ . If instead  $x_8 \in M^1$  then y would be contained in a further cycle of length at most 2D - 1, contradicting Proposition 2.2(i). Therefore,  $M^1$  passes through either  $v_1, v_2$  or  $x_9$ . Consequently, the path  $M^2 = w_2 - x$  goes also through either  $v_1, v_2$ , or  $x_9$ . If both  $M^1$ and  $M^2$  reached x through either  $v_1$  or  $v_2$  then there would exist a cycle of length at most 2D-4 in  $\Gamma$ , a contradiction. We may therefore assume that  $M^1$  goes through  $x_9$ ,  $x_6$  and  $y_1$ , and that  $M^2$  goes through  $v_1$ ; see Fig. 9(a). Note that  $D^3 = w_1 M^1 x_9 x_6 y_1 x_1 v \dots v_1 M^2 w_2 w w_1$  is a 2D-cycle, and that after the above developments, the vertices v and

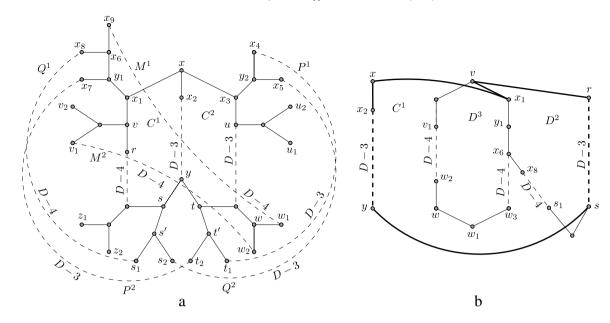
 $x_1$  cannot be contained in any further short cycle, because they are already contained in the cycles  $C^1$ ,  $D^2$  and  $D^3$ .

Let us now turn our attention to the vertex v and the cycles  $C^1$ ,  $D^2$  and  $D^3$ ; see Fig. 9(b) (cycle  $C^1$  is highlighted by a heavier line).

Let  $w_3$  be the neighbor of  $w_1$  different from w such that  $w_3 \in M^1$ . See Fig. 9(b).

We see that the premises of the Intersection Lemma hold. Mapping the vertex v to  $\alpha$ , r to  $\alpha$ , w to  $\beta$ , and mapping the 2D-cycle  $D^3$  to  $\mathcal{D}^1$ , and the (2D-1)-cycle  $C^1$  to  $\mathcal{C}$ , we obtain, by the Intersection Lemma, that one of the following cases holds. In the first case, v and r are contained in a (2D-1)-cycle that intersects with  $D^3$  at a path of length D-1. This cycle would be precisely  $C^1$ , implying D-1=1, a contradiction. In the second case, v and r are contained in another 2D-cycle that intersects with  $D^3$  at a path of length D-1. This cycle would be precisely  $D^2$ , implying D-1=3, a contradiction.

As a result, when  $D \ge 5$ , a (3, D, -4)-graph does not contain a vertex of Type (i).



**Fig. 9.** Auxiliary figure for Proposition 3.1.

Non-existence of vertices of Type (ii)

Proposition 3.1 opens up a way to prove the non-existence of vertices of Type (ii), as shown below.

**Proposition 3.2.** A(3, D, -4)-graph,  $D \ge 5$ , does not contain a vertex of Type (ii).

**Proof.** Let x be a vertex of Type (ii). Let H be a subgraph of  $\Gamma$  isomorphic to  $\Theta_D$ , where x and y are its branch vertices. Then H consists of three independent paths x-y of length D, say  $P^1$ ,  $P^2$  and  $P^3$ . Since x is of Type (ii), x is also contained in one (2D-1)-cycle, say C.

We may assume that  $|V(P^1 \cap C)| > 1$ ,  $|V(P^3 \cap C)| > 1$  and  $V(P^2 \cap C) = \{x\}$ . As C is a (2D-1)-cycle, there is a vertex u of  $P^1$ , different from x or y, such that u and the neighbor of u not contained in  $P^1$ , say  $u_1$ , are both contained in C. Let v and w be the vertices in  $P^2$  and  $P^3$ , respectively, at distance D from u in H. If the distance in  $\Gamma$  between u and either v or w was at most D-1 then u would be contained in two cycles of length at most (2D-1) and in two 2D-cycles, contradicting Proposition 2.2(iii). Therefore, the distance in  $\Gamma$  between u, and v or w is v. Let v and v and v and v and v and v and v or v is v. Let v and v

A path  $P=u_1-v$  does not pass through u, otherwise some vertices of H would be contained in a cycle of length at most 2D-3. Suppose that P passes through  $v_3$ . If  $x\sim v$  (implying  $x=v_3=w_3$ ) then either there would be a cycle of length at most 2D-3 or the distance in  $\Gamma$  between u and w would be at most D-1, a contradiction. If instead  $x\sim v$  then x and some vertices of  $P^2$  would be contained in a cycle of length at most 2D-1, but, by Proposition 2.2(ii), this cycle would be C, contradicting our assumption that  $V(P^2\cap C)=\{x\}$ . Therefore,  $v_3\not\in P$ . If P passed through  $v_2$  then u would be contained in a cycle of length at most 2D-1 that does not contain x. This is a contradiction, because, since u is a vertex of Type (iii), u can be contained in only one (2D-1)-cycle, and that (2D-1)-cycle is C. Therefore, P passes through the neighbor of v not contained in  $P_2$ , say  $v_1$ , and P is a D-path.

Let r be the neighbor of  $u_1$ , different from u, and not contained in P.

Reasoning as above, a path Q=r-v does not contain  $v_2$  or  $v_3$ , otherwise u would be contained in a further short cycle, contradicting Proposition 2.2(iii). As a result, Q uses the edge  $vv_1$ , and is a D-path. Thus,  $u_1$  is contained in an additional (2D-1)-cycle, namely,  $C^1=u_1Pv_1Qru_1$ , and  $u_1$  is therefore a vertex of Type (i) ( $u\in C\cap C^1$ ). However, as  $D\geq 5$ , by Proposition 3.1,  $\Gamma$  does not contain such a vertex; see Fig. 10(b).

Non-existence of vertices of Type (iii)

At this point, assuming that there are no vertices of Type (i) or (ii), our aim is to rule out the existence of vertices of Type (iii).

**Lemma 3.3.** Let C be a (2D-1)-cycle in  $\Gamma$ . Let  $D^1$  be a 2D-cycle such that C and  $D^1$  are non-disjoint. Then the intersection of C and  $D^1$  is a path of length at most D-2.

**Proof.** Let C be a (2D-1)-cycle in  $\Gamma$ . Suppose, by way of contradiction, that in  $\Gamma$  there exists a 2D-cycle, say  $D^1$ , such that the intersection of C and  $D^1$  is a path of length exactly D-1.

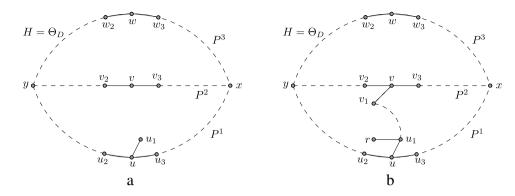


Fig. 10. Auxiliary figure for Proposition 3.2.

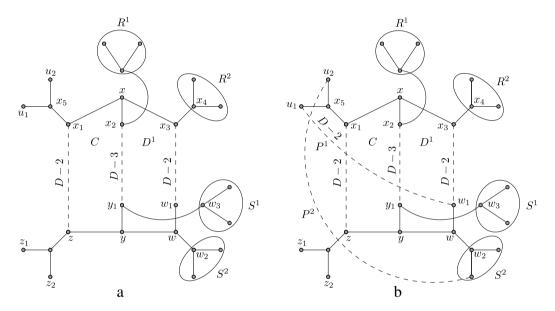


Fig. 11. Auxiliary figure for Lemma 3.3.

Since  $D \ge 5$ ,  $\Gamma$  contains the subgraph depicted in Fig. 11(a). Let x be a vertex lying on C and  $D^1$  such that x has a neighbor, say  $x_3$ , belonging to  $D^1 - C$ , and let y and z be the vertices at distance D - 1 from x in C. Let w be the vertex at distance D from x in  $D^1$ .

Let the vertices  $u_1$ ,  $u_2$ ,  $x_1$ ,  $x_2$ ,  $x_4$ ,  $x_5$ ,  $z_1$ ,  $z_2$ ,  $y_1$ ,  $w_1$ ,  $w_2$  and  $w_3$ , and the sets  $S^1$ ,  $S^2$ ,  $R^1$ , and  $R^2$  be as in Fig. 11(a).

We first consider a path  $P^1 = u_1 - y$ . Note that  $P^1$  cannot pass through  $x_5$ . If  $P^1$  passed through z then  $x_1$  would be contained in another cycle of length at most 2D - 1, a contradiction. Therefore,  $P^1$  passes through either  $w_1$ , a vertex from the set  $S^1$ , or a vertex from the set  $S^2$ .

Suppose that  $w_1 \in P^1$ . In this case  $P^1$  must be a D-path, otherwise x would be contained in a further cycle of length at most 2D - 1, a contradiction. Then x is contained in an additional 2D-cycle, namely,  $D^2 = u_1x_5x_1xx_3D^1w_1P^1u_1$ . Note that x is saturated.

A path  $P^2 = u_2 - y$  does not contain  $y_1$  or  $w_1$ , otherwise in  $\Gamma$  there would be a cycle of length at most 2D - 2 or x would be contained in a further cycle of length at most 2D, a contradiction. Therefore,  $P^2$  is a path of length D - 1 or D which goes through  $S^2$ , and forms the cycle  $D^3 = u_1 P^1 w_1 w w_2$  (a vertex in  $S^2$ )  $P^2 u_2 x_5 u_1$ . The cycle  $D^3$  is a (2D - 1)-cycle or a 2D-cycle, depending on the length of  $P^2$ . Note that, since  $w_1 \in D^1$ ,  $D^2$ ,  $D^3$ , the vertex  $w_1$  cannot be contained in any further (2D - 1)-cycle. See Fig. 11(b).

In this case a path  $x_4 - y$  cannot go through w (that would form a new cycle of length at most 2D - 1 containing  $w_1$ ), neither can  $x_4 - y$  pass through  $y_1$  or z (otherwise x would be contained in a further cycle of length at most 2D). Consequently,  $d(x_4, y) > D$ , a contradiction.

As a result,  $w_1 \notin P^1$ , and  $P^1$  reaches y through either a vertex from  $S^1$  or a vertex from  $S^2$ . We can then assume that  $P^2 = u_2 - y$  also reaches y through either a vertex from  $S^1$  or a vertex from  $S^2$ . Note that  $P^1$  and  $P^2$  intersect neither in  $S^1$  nor in  $S^2$ , otherwise there would exist a cycle of length at most 2D - 2. We may accordingly assume that  $P^1$  goes through a

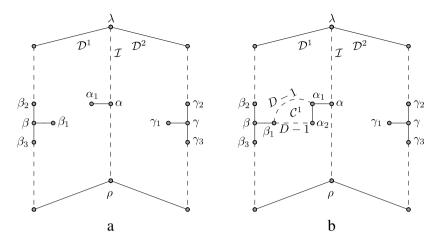


Fig. 12. Auxiliary figure for Saturation Lemma.

vertex in  $S^1$ , and that  $P^2$  goes through a vertex in  $S^2$ . Then,  $P^1$  must be a D-path, whereas  $P^2$  could be either a (D-1)-path or a D-path.

In this way, we have obtained a new 2*D*-cycle containing *x* and  $x_2$ , say,  $D^2 = u_1 P^1$  (a vertex in  $S^1$ )  $w_3 y_1 C x_2 x x_1 x_5 u_1$ . Note that the vertices *x* and  $x_2$  are both saturated.

Finally, we consider the paths  $Q^1 = z_1 - x$  and  $Q^2 = z_2 - x$ .

Reasoning as before, the paths  $Q^1$  and  $Q^2$  reach x through either a vertex from the set  $R^1$  or a vertex from the set  $R^2$ , but both paths cannot go through the same set. Therefore, we may assume that  $Q^1$  passes through a vertex in  $R^1$ , and that  $Q^2$  goes through a vertex in  $R^2$ . But in this case,  $x_2$  would be contained in a further cycle of length at most 2D, a contradiction.

Thus, the lemma follows.

Next we prove a lemma that will be very useful from now on.

**Lemma 3.4** (Saturation Lemma). Let  $\mathcal{D}^1$  and  $\mathcal{D}^2$  be two 2D-cycles intersecting at a path 1 of length D-1. Let  $\lambda$  and  $\rho$  be the vertices lying on 1 at distance D-1 from each other. Suppose that  $\lambda$  is saturated and that there exists a vertex  $\alpha \neq \lambda$ ,  $\rho$  lying on 1 such that its neighbor  $\alpha_1$  not contained in 1 does not belong to any of the short cycles saturating  $\lambda$ . Then the following two assertions hold:

- (i) There is at least one further short cycle  $\mathcal{D}^3$  containing  $\alpha$ ,  $\alpha_1$  and  $\rho$ .
- (ii) If  $\eta_1$  is the neighbor not contained in 1 of a vertex  $\eta \in 1$  such that  $\eta \neq \lambda, \alpha, \rho$ , then  $\eta_1$  does not belong to  $\mathcal{D}^3$ .

**Proof.** Let  $\mathcal{D}^1$  and  $\mathcal{D}^2$  be two 2*D*-cycles intersecting at a path  $\mathcal{I}$  of length D-1, and let  $\lambda$  and  $\rho$  be the vertices lying on  $\mathcal{I}$  at distance D-1 from each other. Suppose that  $\lambda$  is saturated, and that there exists a vertex  $\alpha \neq \lambda$ ,  $\rho$  lying on  $\mathcal{I}$  such that its neighbor  $\alpha_1$  not contained in  $\mathcal{I}$  does not belong to any of the short cycles saturating  $\lambda$ .

Let  $\beta$  and  $\gamma$  be the vertices in  $\mathcal{D}^1$  and  $\mathcal{D}^2$ , respectively, at distance D from  $\alpha$ , and let the vertices  $\beta_1$ ,  $\beta_2$ ,  $\beta_3$ ,  $\gamma_1$ ,  $\gamma_2$  and  $\gamma_3$  be as in Fig. 12(a).

Suppose, on the contrary, that there is no short cycle containing  $\alpha$ ,  $\alpha_1$  and  $\rho$ .

Consider a path  $P^1 = \alpha_1 - \beta$ . Note that  $\alpha \notin P^1$ . If  $P^1$  went through  $\beta_2$  then  $\lambda$  would be contained in a further short cycle, contradicting the saturation of  $\lambda$ . If instead  $P^1$  passed through  $\beta_3$  then  $\rho$  would belong to a short cycle that also contains  $\alpha$  and  $\alpha_1$ , a contradiction. Therefore,  $P^1$  reaches  $\beta$  through  $\beta_1$ , and is D-path.

Let  $\alpha_2$  be the neighbor of  $\alpha_1$ , other than  $\alpha$ , which is not contained in  $P^1$ . Consider a path  $P^2 = \alpha_2 - \beta$ . Note that  $\alpha_1 \notin P^2$ . Then  $P^2$  does not go through  $\beta_2$ , otherwise  $\lambda$  would be contained in a further short cycle. Neither does  $P^2$  pass through  $\beta_3$ , otherwise  $\rho$  would belong to a short cycle that also contains  $\alpha$  and  $\alpha_1$ . Therefore,  $P^2$  reaches  $\beta$  through  $\beta_1$ , and is D-path. Note that  $P^2$  causes the formation of a (2D-1)-cycle  $C^1 = \alpha_1 P^1 \beta_1 P^2 \alpha_2 \alpha_1$ ; see Fig. 12(b).

Following the same analysis as in the case of the paths  $P^1$  and  $P^2$ , we obtain that D-paths  $Q^1 = \alpha_1 - \gamma$  and  $Q^2 = \alpha' - \gamma$  reach  $\gamma$  through  $\gamma_1$ , where  $\alpha'$  is the neighbor of  $\alpha_1$ , other than  $\alpha$ , which is not contained in  $Q^1$ . Consequently, we obtain a new (2D-1)-cycle  $C^2 = \alpha_1 Q^1 \gamma_1 Q^2 \alpha' \alpha_1$ , and thus,  $\alpha_1$  is contained in two (2D-1)-cycles  $C^1$  and  $C^2$ , contradicting Proposition 3.1.

Thus, there is at least one further short cycle containing  $\alpha$ ,  $\alpha_1$  and  $\rho$ , and (i) follows.

Note that the second assertion follows immediately from the proof of (i).

We are now in a position to rule out the existence of vertices of Type (iii).

**Proposition 3.3.** A(3, D, -4)-graph, D > 5, does not contain a vertex of Type (iii).

**Proof.** Let x be a vertex of Type (iii) lying on a (2D-1)-cycle C and two 2D-cycles  $D^1$  and  $D^2$ . Let  $x_1$  be the neighbor of x such that  $x_1 \in (C-D^1)$ , and  $y_1$  the vertex on  $D^1$  at distance D from x.

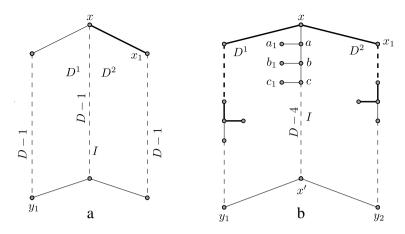


Fig. 13. Auxiliary figure for Proposition 3.3. Parts belonging to the cycle C are highlighted by a heavier line.

To prove the proposition we have prepared the following two claims.

**Claim 1.** The intersection of  $D^1$  and  $D^2$  is a path of length D-1.

**Proof of Claim 1.** We apply the Intersection Lemma. Mapping the vertex x to  $\alpha$ ,  $x_1$  to  $\alpha_1$ ,  $y_1$  to  $\beta$ , and mapping the 2D-cycle  $D^1$  to  $\mathcal{D}^1$ , and the (2D-1)-cycle C to C, we see, by the Intersection Lemma (ii), that the 2D-cycle  $D^2$  intersects  $D^1$  at a path C of length C 1. Note that the case (i) of the Intersection Lemma does not hold because of Lemma 3.3. See Fig. 13(a) (the edge C is highlighted by a heavier line).

**Claim 2.**  $C \cap I = \{x\}.$ 

**Proof of Claim 2.** We use the Intersection Lemma again. Suppose, on the contrary, that  $|V(C \cap I)| > 1$ . Then there are two vertices  $z \neq x$  and  $z_1$  such that  $z \in (C \cap I)$ ,  $z_1 \in (C - I)$  and  $z \sim z_1$ . This implies that z is a vertex of Type (iii), which belongs to C,  $D^1$  and  $D^2$ . Therefore, z is saturated. In this case the premises of the Intersection Lemma hold again. Mapping the vertex z to  $\alpha$ ,  $z_1$  to  $\alpha_1$ , the vertex in  $D^1$  at distance D from z to  $\beta$ , and mapping the (2D-1)-cycle C to C and the D-cycle  $D^1$  to  $D^1$ , we obtain, by the Intersection Lemma (ii), that there exists an additional D-cycle containing D and D are contradiction. This completes the proof of the Claim.

Let x' be the vertex on I at distance D-1 from x, a the neighbor of x contained in I, and  $a_1$  the neighbor of a not contained in I. Since x is saturated, we see that the premises of the Saturation Lemma hold. Mapping the vertex x to  $\lambda$ , x' to  $\rho$ , a to  $\alpha$  and  $a_1$  to  $\alpha_1$ , and mapping the 2D-cycle  $D^1$  to  $\mathcal{D}^1$  and  $D^2$  to  $\mathcal{D}^2$ , it follows that there is a further short cycle  $D^3$  containing a,  $a_1$  and x'.

Let  $b \neq x$  be the neighbor of a contained in I,  $b_1$  the neighbor of b not contained in I,  $c \neq a$  the neighbor of b contained in I and  $c_1$  the neighbor of c not contained in I. Since  $D \geq 5$ , it follows that  $a \neq x, x', b \neq x, x'$  and  $c \neq x, x'$ . See Fig. 13(b).

By Saturation Lemma (ii), we see that neither  $b_1$  nor  $c_1$  is contained in  $D^3$ . Therefore, we can apply the Saturation Lemma again. Mapping the vertex x to  $\lambda$ , x' to  $\rho$ , b to  $\alpha$  and  $b_1$  to  $\alpha_1$ , and mapping the 2D-cycle  $D^1$  to  $D^1$  and  $D^2$  to  $D^2$ , it follows that there is a further short cycle  $D^4$  containing b,  $b_1$  and x'. Therefore, x' is saturated.

By Saturation Lemma (ii), we see that  $c_1$  is not contained in  $D^4$ , allowing a further application of the Saturation Lemma. Mapping the vertex x to  $\lambda$ , x' to  $\rho$ , c to  $\alpha$  and  $c_1$  to  $\alpha_1$ , and mapping the 2D-cycle  $D^1$  to  $D^1$  and  $D^2$  to  $D^2$ , it follows that there is a further short cycle  $D^5$  containing c,  $c_1$  and  $c_2$  and  $c_3$  but the formation of the cycle  $c_3$  contradicts the fact that  $c_4$  is saturated.

Thus,  $\Gamma$  does not contain a vertex of Type (iii), and the proposition follows.

We are now in a position to prove the main result of this section.

**Proof of Theorem 3.1.** Combining Propositions 3.1–3.3, the theorem follows.

As an immediate corollary of Theorem 3.1, we obtain that if a (3, D, -4)-graph with  $D \ge 5$  exists then Conjecture 2.1 is true.

**Corollary 3.1.** If a (3, D, -4)-graph with D > 5 exists then it must have girth 2D.  $\square$ 

#### 4. Non-existence of (3, D, -4)-graphs for $D \ge 5$

From Theorem 3.1, it follows that  $\Gamma$  contains only vertices of Type (iv), (v) or (vi). By ruling out the existence of such vertices, we obtain the non-existence of (3, D, -4)-graphs for  $D \ge 5$ .

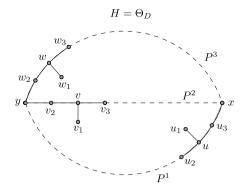


Fig. 14. Auxiliary figure for Proposition 4.1.

**Proposition 4.1.** A(3, D, -4)-graph,  $D \ge 5$ , does not contain a subgraph isomorphic to  $\Theta_D$ .

**Proof.** Let *H* be a subgraph of  $\Gamma$  isomorphic to  $\Theta_D$ , where *x* and *y* are its branch vertices.

Let u be the vertex on  $P^1$  at distance 2 from x. Let v and w be vertices on  $P^2$  and  $P^3$ , respectively, such that d(u, v) = d(u, w) = D. Let  $u_2$  and  $u_3$ ,  $v_2$  and  $v_3$ , and  $w_2$  and  $w_3$  be the neighbors of u on  $P^3$ , the neighbors of v on v0, and the neighbors of v0 on v1, respectively. Let v1, v1 and v3 be the neighbors of v3, respectively, that do not belong to v3, see Fig. 14.

First, consider a path  $Q^1 = u_1 - v$ . Then,  $Q^1$  does not go through u,  $v_2$  or  $v_3$ , otherwise there would exist a cycle of length at most 2D - 1 in  $\Gamma$ . Therefore,  $Q^1$  goes through  $v_1$ . Suppose that  $Q^1$  is a D-path.

Let  $r \neq u$  and  $s \neq v$  be the neighbors of  $u_1$  and  $v_1$ , respectively, that do not belong to  $Q^1$ .

A path  $Q^2 = r - v$  does not pass through  $v_1$ , otherwise there would exist a cycle of length at most 2D - 1 in  $\Gamma$ . Then  $Q^2$  passes through either  $v_2$  or  $v_3$ , and should be a D-path, otherwise there would exist a cycle of length at most 2D - 1 in  $\Gamma$ . Analogously, a path  $Q^3 = s - u$  is a D-path, and goes through either  $u_2$  or  $u_3$ .

Note that the paths  $Q^2$  and  $Q^3$  form part of two 2D-cycles, denoted by  $D^1$  and  $D^2$ , which contain either x or y. The cycle  $D^1$  is either  $uu_3P^1xP^2v_3Q^2ru_1u$  or  $uu_2P^1yP^2v_2Q^2ru_1u$ , while the cycle  $D^2$  is either  $u_2P^1yP^2v_2vv_1sQ^3u_2$  or  $u_3P^1xP^2v_3vv_1sQ^3u_3$ . Note that the cycles  $D^1$  and  $D^2$  do not contain  $w_1$ ,  $w_2$  or  $w_3$ .

Let us further suppose that a path  $T^1 = u_1 - w$  is a D-path. By following the same reasoning as in the case of the paths  $Q^1$ ,  $Q^2$  and  $Q^3$ , we obtain that  $T^1$  passes through  $w_1$ , and that there are two further 2D-cycles, say,  $D_3$  and  $D_4$ , containing either x or y. In this case the vertices x and y are of Type (v), and, by Proposition 2.2(v), x and y are saturated.

Since  $D \ge 5$ , we can find another vertex on  $P^1$ , say z, different from x,  $u_3$ , u,  $u_2$  or y. Let p be the vertex on  $P^2$  such that d(z, p) = D, and  $z_1$  the neighbor of z that does not belong to H. Note that  $z_1$  does not belong to  $D^1$ ,  $D^2$ ,  $D^3$  or  $D^4$ .

We consider paths  $R^1 = z_1 - p$  and  $R^2 = q - p$ , where  $q \neq z$  is the neighbor of  $z_1$  not contained in  $R^1$ . Note that the paths  $R^1$  and  $R^2$  must be D-paths, otherwise x and y would be contained in a further short cycle. Then we obtain a new 2D-cycle containing  $z_1$  and either x or y, a contradiction to Proposition 2.2(v). Therefore,  $T^1$  is a (D-1)-path.

If the path  $T^1$  is a (D-1)-path then there are two new 2D-cycles containing  $w_1$ , different from  $D^1$  or  $D^2$ , such that one contains x, and the other contains y. Therefore, as before, vertices x and y are both of Type (v), and are saturated. We consider again the aforementioned vertices z,  $z_1$ , p and q, and the paths  $R^1 = z_1 - p$  and q - p. In this case  $z_1$  does not belong to any of the cycles involving x or y. As a result, we obtain a new 2D-cycle containing  $z_1$  and either x or y, a contradiction.

Thus,  $Q^1$  is a (D-1)-path and so is  $T^1$ .

By analogy, if the paths  $Q^1$  and  $T^1$  are (D-1)-paths then there are four new 2D-cycles such that two of them contain x, and the other two contain y. Therefore, x and y cannot be contained in any additional short cycle. However, we can again use the vertices z,  $z_1$ , p and q, and the paths  $R^1 = z_1 - p$  and q - p to find a further 2D-cycle containing  $z_1$  and either x or y, contradicting Proposition 2.2(v).

**Corollary 4.1.** A (3, D, -4)-graph, D > 5, does not contain a vertex of Type (iv)or (v).  $\Box$ 

**Proposition 4.2.** A(3, D, -4)-graph, D > 5, does not contain a vertex of Type (vi).

**Proof.** Let x be a vertex of  $\Gamma$ . Then x is a vertex of Type (vi). Let  $D^1$  be one of the 2D-cycles on which x lies, and  $y_1$  the vertex in  $D^1$  at distance D from x. Furthermore, we denote by  $w_2$  the neighbor of x not contained in  $D^1$ .

In this case, by the Intersection Lemma, mapping the vertex x to  $\alpha$ ,  $w_2$  to  $\alpha_1$ ,  $y_1$  to  $\beta$ , and mapping the 2D-cycle  $D^1$  to  $\mathcal{D}^1$  ( $w_2$  belongs to no (2D-1)-cycle), we see that there exists another 2D-cycle containing x and  $w_2$ , say  $D^2$ , such that the intersection of  $D^1$  and  $D^2$  is a path of length D-1.

We prove this proposition by reasoning in the same way as in the proof of Proposition 4.1.

Let u and w be the vertices in  $D^1 - D^2$  and in  $D^2 - D^1$ , respectively, at distance 2 from x. Let v be the vertex in  $D^1 \cap D^2$  such that d(u, v) = d(w, v) = D. Let  $v_3$  be the vertex in  $D^1 \cap D^2$  at distance D - 1 from x. Finally, let the vertices  $u_1, u_2, u_3, v_1, v_2, w_1, w_3$  and  $y_2$  be as in Fig. 15.

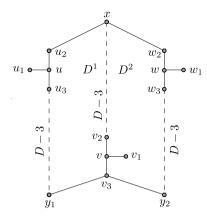


Fig. 15. Auxiliary figure for Proposition 4.2.

Consider a path  $P^1 = u_1 - v$ . Then  $P^1$  does not go through u,  $v_2$  or  $v_3$ , otherwise there would exist a cycle of length at most 2D - 1 in  $\Gamma$ . Therefore,  $P_1$  goes through  $v_1$ . If  $P^1$  was a (D - 1)-path then both u and v would be branch vertices of a  $\Theta_D$ , contradicting Proposition 4.1. As a result,  $P^1$  is a D-path.

Let  $r \neq u$  and  $s \neq v$  be the neighbors of  $u_1$  and  $v_1$ , respectively, that do not belong to  $P^1$ .

A path  $P^2 = r - v$  does not pass through  $v_1$ , otherwise there would exist a cycle of length at most 2D - 1 in  $\Gamma$ . Then  $P^2$  intersects  $D^1$  at v and either  $v_2$  or  $v_3$ , and should be a D-path, otherwise there would exist a cycle of length at most 2D - 1 in  $\Gamma$ . Analogously, a path  $P^3 = s - u$  is a D-path, and intersects  $D^1$  at u and either  $u_2$  or  $u_3$ .

Note that the paths  $P^2$  and  $P^3$  form part of two 2D-cycles, denoted by  $D^3$  and  $D^4$ , which contain either x or  $v_3$ . The cycle  $D^3$  is either  $uu_3D^1v_1v_3v_2P^2ru_1u$  or  $uu_2xD^1v_2P^2ru_1u$ , while the cycle  $D^4$  is either  $uu_3D^1v_1v_3v_2P^3u_2$  or  $uu_3D^1v_1v_3v_1sP^3u_3$ .

Note that the cycles  $D^3$  and  $D^4$  do not contain  $w_1$ ,  $w_2$  or  $w_3$ .

Consider a path  $T^1 = w_1 - v$ . By following the same reasoning as in the case of the paths  $P^1$ ,  $P^3$  and  $P^3$ , we obtain that  $T^1$  passes through  $v_1$ , that  $T^1$  is a D-path and that there are two further 2D-cycles, say,  $D_5$  and  $D_6$ , containing either x or  $v_3$ . In this case the vertices x and  $v_3$  are saturated.

Since  $D \ge 5$ , we can find another vertex in  $D^1 - D^2$ , say z, different from  $u_2$ , u,  $u_3$  or  $y_1$ . Let p be the vertex in  $D^1 \cap D^2$  such that d(z, p) = D, and  $z_1$  the neighbor of z that does not belong to  $D^1$ . Note that  $z_1$  does not belong to  $D^3$ ,  $D^4$ ,  $D^5$  or  $D^6$ .

Then, by considering the paths  $R^1 = z_1 - p$  and q - p, where  $q \neq z$  is the neighbor of  $z_1$  not contained in  $R^1$ , we obtain a new 2*D*-cycle containing  $z_1$  and either x or  $v_3$ , a contradiction to Proposition 2.2(vi).

Thus in a (3, D, -4)-graph with D > 5 there exists no vertex of Type (vi), and the proposition follows.

Combining the results of Theorem 3.1, Corollary 4.1 and Proposition 4.2, we obtain the main result of this paper (Theorem 4.1), thus completing the catalogue of (3, D, -4)-graphs with  $D \ge 2$ .

**Theorem 4.1.** For  $D \ge 5$  there are no (3, D, -4)-graphs.

## 5. Conclusions

In this paper, by proving the non-existence of (3, D, -4)-graphs with  $D \ge 5$ , we have completed the census of (3, D, -4)-graphs with  $D \ge 2$  and  $\epsilon \le 4$ , which is summarized below.

**Catalogue of** (3, D, 0)-**graphs with**  $D \ge 2$ . With the exceptions of the complete graph on 4 vertices and the Petersen graph, there is no cubic Moore graph.

**Catalogue of** (3, D, -2)-**graphs with**  $D \ge 2$ . There are only three non-isomorphic (3, D, -2)-graphs with  $D \ge 2$ ; all shown in Fig. 1.

**Catalogue of** (3, D, -4)-**graphs with**  $D \ge 2$ . For diameter 2 there exist two regular (graphs (a) and (b) in Fig. 2) and three non-regular (3, 2, -4)-graphs (graphs (c), (d) and (e) in Fig. 2). When the diameter is 3, there is a unique (3, 3, -4)-graph; see Fig. 3. The results of this paper, combined with [11], assert that there are no (3, D, -4)-graphs with  $D \ge 4$ .

Contribution to the degree/diameter problem

Our result also improves the upper bound on  $N_{3,D}$ ,  $D \ge 5$ , implying that any maximal graph of maximum degree 3 and diameter  $D \ge 5$  must have order at most  $M_{3,D} - 6$ .

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